

Cognitive Science Knowledge Representation & Organisation

Erasmus Programme
Professor Markos Dendrinou

Predicate Logic

Universe of discourse and typical language

A predicate is a relation with one or more parameters corresponding to a fact:

blue(sky) : the sky is blue

sky is a constant

csg (C,S,G) : Course-Student-Grade. Student with the identity S has taken in the course C grade G. For certain values of C,S,G (replacing variables with constants) it is true or false

snap (S, N, A, P) : Student-Name-Address-Phone. Student with identity S has name N, address A and phone number P. For certain values of S, N, A, P (replacing variables with constants) it is true or false

human(X) : X is human

mortal (X) : X is mortal

C, S, G, S, N, A, P, X are variables which can be substituted by constants

Quantifiers

Universal quantifier

$\forall X (\text{man}(X) \rightarrow \text{mortal}(X))$

Existential quantifier

$\exists X (\text{year}(X) \rightarrow \text{leap}(X))$

Deduction Rules

1. Application of Modus Ponens

$\text{human}(a) \rightarrow \text{mammal}(a)$

$\text{human}(a)$

We conclude:

$\text{mammal}(a)$

2. Incorrect application of Modus Ponens

$\text{human}(a) \rightarrow \text{mammal}(a)$

$\text{mammal}(a)$

We conclude:

$\text{human}(a)$ [NOT VALID!!!]

3. Application of Modus Tollens

$\text{human}(a) \rightarrow \text{mammal}(a)$

NOT $\text{mammal}(a)$

We conclude (M.T.):

NOT $\text{human}(a)$

Substitution & M.P.

Given propositions:

If something (X) is human then this is mortal

$(\forall X) (\text{human}(X) \rightarrow \text{mortal}(X))$ (prop.1)

Vincent is human

$\text{human}(\text{vincent})$ (prop.2)

To be proved:

Vincent is mortal

$\text{mortal}(\text{vincent})$ (prop.3)

(prop.1) substitute X with vincent:

$\text{human}(\text{vincent}) \rightarrow \text{mortal}(\text{vincent})$ (prop.4)

From (prop.4) and (prop.2) through M.P. we conclude that:

$\text{mortal}(\text{vincent})$

Substitution & M.P. (cont)

Given:

Every human is animal

Every greek is human

Markos is greek

To be proved:

Markos is animal

Written in Predicate Logic

$(\forall X) (h(X) \rightarrow a(X))$ (prop.1)

$(\forall X) (g(X) \rightarrow h(X))$ (prop.2)

$g(\text{markos})$ (prop.3)

$a(\text{markos})$

Proof

(prop.2) substitute X with markos

$g(\text{markos}) \rightarrow h(\text{markos})$ (prop.4)

From (prop.4) and (prop.3) through M.P. we conclude that:

$h(\text{markos})$ (prop.5)

(prop.1) substitute X with markos

$h(\text{markos}) \rightarrow a(\text{markos})$ (prop.6)

From (prop.6) and (prop.5) through M.P. we conclude that:

$a(\text{markos})$

Exercise-3a

Given propositions:

If something (X) is cat then this is feline

Eleanor is cat

To be proved:

Eleanor is feline

The proof must be presented by using the deduction rule: Modus Ponens (M.P.) along with substitution

Exercise-3b

- Someone who gets high degree in medicine becomes good physician
- A good physician cures the patients
- Anne got high degree in medicine
- To be proved
- Anne cures the patients

De Morgan rules

- $\neg (p \vee q) \models \neg p \wedge \neg q$
- $\neg (p \wedge q) \models \neg p \vee \neg q$
- $\neg(\forall X) (p(X)) \models (\exists X) (\neg p(X))$
- $\neg(\exists X) (p(X)) \models (\forall X) (\neg p(X))$

Resolution principle: Elimination of quantifiers

- We can eliminate universal quantifier by substituting the variable name with a new name, uniquely defined for each proposition.
- We can eliminate existential quantifier by substituting the variable name with an assumed constant corresponding to a certain object of the universe of discourse, uniquely defined for each proposition.

Proof in Predicate Calculus through Resolution Principle

- Aristotelian scheme
- Given propositions:
 - Every human is animal
 - Every Greek is human
- To be proved:
 - Every Greek is animal
- Step 1
 - $(\forall X) (h(X) \rightarrow an(X))$ (p1)
 - $(\forall X) (g(X) \rightarrow h(X))$ (p2)
 - $\neg ((\forall X) (g(X) \rightarrow an(X)))$ (p3)

Proof in Predicate Calculus through Resolution Principle (cont)

- Step 2

- $(\forall X) (\neg h(X) \vee an(X))$ (p1)
- $(\forall X) (\neg g(X) \vee h(X))$ (p2)
- $(\exists X) \neg (g(X) \rightarrow an(X)) \models$
- $(\exists X) \neg (\neg g(X) \vee an(X)) \models$
- $(\exists X) (g(X) \wedge \neg an(X))$

Proof in Predicate Calculus through Resolution Principle (cont)

● Step 3

- $\neg h(X1) \vee an(X1)$ (p1)
- $\neg g(X2) \vee h(X2)$ (p2)
- $g(a) \wedge \neg an(a) \models$
- $g(a)$ (p3)
- $\neg an(a)$ (p4)

● Step 4

- (p1) and (p4) for $X1=a \Rightarrow \neg h(a)$ (p5)
- (p2) and (p3) for $X2=a \Rightarrow h(a)$ (p6)
- (p5) and (p6) $\Rightarrow 0$ (αντίφαση)

Proof through Resolution Principle

–Exercise 4

- If every M are P and some S are M \Rightarrow some S are P”.
- Each m.p. is of high salary
- $(\forall X) (mp(X) \rightarrow hs(X))$ (prop1)
- Some lawyers are m.p
- $(\exists X) (l(X) \wedge mp(X))$ (prop2)
- Conclusion (to be proved)
- Some lawyers are of high salary
- $(\exists X) (l(X) \wedge hs(X))$ (prop3)
-